

DNA evolution, Automata and Clumps

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Problem setting

- ▶ Alphabet $\mathcal{A} = \{\text{A, C, G, T}\}$ (DNA)

time = 0 $S_n(0) =$ **YYYYYYY.....YYYYYYY**

\vdots \vdots \vdots

time = T $S_n(T) =$ **YYYY...FF..FF..YYYYYY**

- ▶ n = length of random sequences $S_n(0), \dots, S_n(T)$ ($n \approx 2000$)
- ▶ $b = \text{FF..FF} \in \mathcal{A}^k$ Transcription Factor ($5 \leq k = |b| \leq 10$)
 - ▶ b **does not occur** in $S_n(0)$
 - ▶ b **occurs for the first time by evolution** at time T in a sequence evolving from $S_n(0)$

Aim: Compute T

Initial $\nu(\alpha)$ and Substitution Probabilities $\pi_{\alpha \rightarrow \beta}$

α	$\nu(\alpha)$
A	0.23889
C	0.26242
G	0.25865
T	0.24004



substitution prob.

$$\mathbb{P}(1) = \pi_{\alpha \rightarrow \beta}$$

for **one** generation
(20 years)

A		A	0.9999999763
A		C	$4.54999994943 \times 10^{-9}$
A		G	$1.57499995613 \times 10^{-8}$
A		T	$3.40000001733 \times 10^{-9}$
C		A	$6.14999993408 \times 10^{-9}$
C		C	0.99999996495
C		G	$7.14999984731 \times 10^{-9}$
C		T	$2.17499993935 \times 10^{-8}$
G		A	$2.17499993935 \times 10^{-8}$
G		C	$7.14999984731 \times 10^{-9}$
G		G	0.99999996495
G		T	$6.14999993408 \times 10^{-9}$
T		A	$3.40000001733 \times 10^{-9}$
T		C	$1.57499995613 \times 10^{-8}$
T		G	$4.54999994943 \times 10^{-9}$
T		T	0.9999999763

Powers of $\mathbb{P}(1)$ remains close to the Identity Matrix

$$\mathbb{P}(1) \approx \begin{pmatrix} 1 - 3m & m & m & m \\ m & 1 - 3m & m & m \\ m & m & m & 1 - 3m \\ m & m & m & 1 - 3m \end{pmatrix} \quad \text{with } m \approx 10^{-8}$$

$$\mathbb{P}^N(1) \approx \begin{pmatrix} 1 - 3mN & mN & mN & mN \\ mN & 1 - 3mN & mN & mN \\ mN & mN & mN & 1 - 3mN \\ mN & mN & mN & 1 - 3mN \end{pmatrix} + \mathcal{O}(m^2N)$$

Therefore

$$P^N(1) \times \nu \approx \nu \quad \text{for } N \approx 10^6 \text{ and } N < 10^6$$

$$P^\infty(1) \times \nu = (0.25, 0.25, 0.25, 0.25)^t$$

Geometric distribution of the Waiting Time

By **stationnarity** of ν , assuming $T \in \mathbb{N}$

$$\begin{aligned} & \mathbf{P}(\text{no } b \text{ in } S_n(j+1) \mid \text{no } b \text{ in } S_n(j)) \\ &= \mathbf{P}(\text{no } b \text{ in } S_n(1) \mid \text{no } b \text{ in } S_n(0)) \\ &= 1 - \mathbf{P}(b \text{ occurs in } S_n(1) \mid \text{no } b \text{ in } S_n(0)) \end{aligned}$$

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$$\begin{aligned} & \mathbf{P}(b \text{ occurs in } S_n(T) \mid \text{no } b \text{ in } S_n(T-1)) \\ &= \mathbf{P}(b \text{ occurs in } S_n(1) \mid \text{no } b \text{ in } S_n(0)) \end{aligned}$$

Geometric distribution of the Waiting Time

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$$\begin{aligned} & \mathbf{P}(b \text{ occurs in } S_n(T) \mid \text{no } b \text{ in } S_n(T-1)) \\ &= \mathbf{P}(b \text{ occurs in } S_n(1) \mid \text{no } b \text{ in } S_n(0)) \end{aligned}$$

Setting $\mathbf{p}_n = \mathbf{P}(b \text{ occurs in } S_n(1) \mid \text{no } b \text{ in } S_n(0))$,

$$\mathbf{E}(T) \approx \sum_{i \geq 0} (1 - \mathbf{p}_n)^i \times \mathbf{p}_n = \frac{1}{\mathbf{p}_n}$$

Renewing the Aim

We need now computing

$$\begin{aligned} p_n &= \mathbf{P}(b \text{ occurs in } S_n(1) \mid \text{no } b \text{ in } S_n(0)) \\ &= \frac{\mathbf{P}(b \text{ occurs in } S_n(1) \text{ AND no } b \text{ in } S_n(0))}{\mathbf{P}(\text{no } b \text{ in } S_n(0))} \end{aligned}$$

Different computations of p_n

1. Behrens-Vingron (2010)
 - ▶ Approach **neglecting words correlation**.
 - ▶ **Efficient computation** of p_n with respect to this assumption.
2. Behrens-Nicaud-N (2012)
 - ▶ **Rigorous and efficient approach by automata**.
3. N (NCMA2012)
 - ▶ Heuristic approach by **clump analysis**, either by **combinatorics of words** or by **automata** and generating functions.
4. N (2013)
 - ▶ Heuristic approach, adaptation of the **Régnier-Szpankowski equations** and **explicit formula** approximating p_n

Different computations of p_n

$$\begin{array}{lcl} \text{time} = 0 & S_n(0) = & \text{YYYYYYYY} \dots \text{YYYYYYYY} \\ & \vdots & \\ & \vdots & \\ \text{time} = 1 & S_n(1) = & \text{YYYY} \dots \text{FF} \dots \text{FF} \dots \text{YYYY} \end{array}$$

- ▶ Behrens-Vingron compute the probability that b occurs in $S_n(1)$ (**without allowing overlaps of occurrences**), and then the probability that $S_n(0)$ evolves to $S_n(1)$

Different computations of p_n

$$\begin{array}{lcl} \text{time} = 0 & S_n(0) = & \text{YYYYYYY} \dots \text{YYYYYYY} \\ & \vdots & \vdots \\ & \vdots & \vdots \\ \text{time} = 1 & S_n(1) = & \text{YYYY} \dots \text{FF} \dots \text{FF} \dots \text{YYYYY} \end{array}$$

- ▶ Behrens-Vingron compute the probability that b occurs in $S_n(1)$ (**without allowing overlaps of occurrences**), and then the probability that $S_n(0)$ evolves to $S_n(1)$
- ▶ Behrens-Nicaud-N use an automaton on the alphabet $\mathcal{A} \times \mathcal{A}$ that scans **simultaneously** $S_n(0)$ and $S_n(1)$. This automaton is a kind of product of two **Knuth-Morris-Pratt automata**.

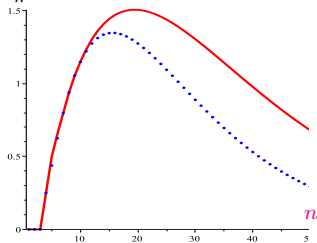
Different computations of p_n

$$\begin{array}{lcl} \text{time} = 0 & S_n(0) = & \text{YYYYYYY} \dots \text{YYYYYYY} \\ & \vdots & \vdots \\ & \vdots & \vdots \\ \text{time} = 1 & S_n(1) = & \text{YYYY} \dots \text{FF} \dots \text{FF} \dots \text{YYYYY} \end{array}$$

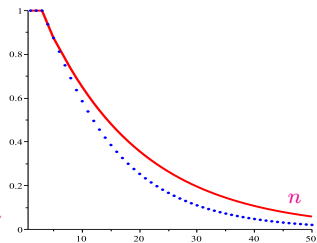
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- ▶ N (2012) assumes that **a single mutation occurred** and considers the **clumps of neighbors** of b at **distance 1** in $S_n(0)$.

An unexpected behaviour

$\frac{p_n}{\pi} \times \mathbf{P}(\text{no } b \text{ (or } b') \text{ in } S_n(0))$



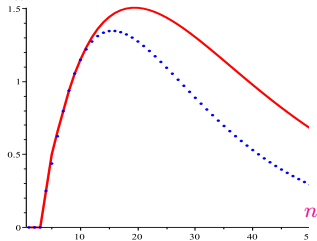
$\mathbf{P}(\text{no } b \text{ (or } b') \text{ in } S_n(0))$



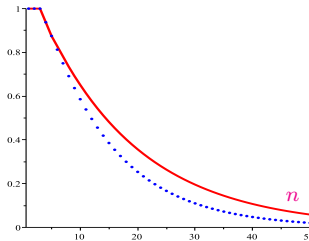
$$b = \text{ACAC} \quad b' = \text{AACC}, \quad \nu(\text{A}) = \nu(\text{C}) = \frac{1}{2}, \quad \pi = \pi_{\text{A} \rightarrow \text{C}} = \pi_{\text{C} \rightarrow \text{A}}$$

An unexpected behaviour

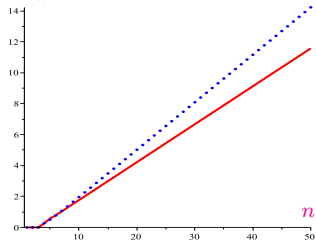
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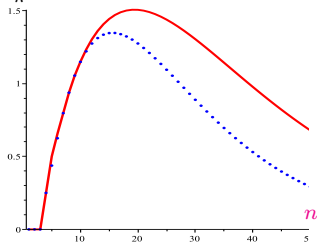
$\frac{p_n}{\pi}$



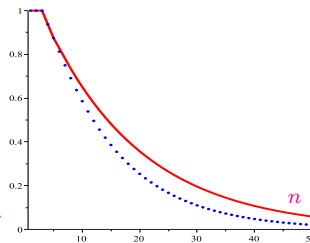
$$b = ACAC \quad b' = AACC, \quad \nu(A) = \nu(C) = \frac{1}{2}, \quad \pi = \pi_{A \rightarrow C} = \pi_{C \rightarrow A}$$

An unexpected behaviour

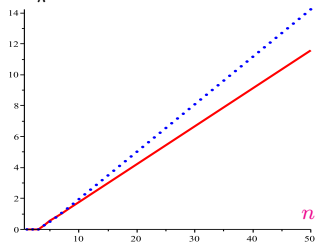
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$\mathbf{P}(\text{no } b \text{ (or } b') \text{ in } S_n(0))$



$\frac{p_n}{\pi}$



$$b = \text{ACAC} \quad b' = \text{AACC}, \quad \nu(\text{A}) = \nu(\text{C}) = \frac{1}{2}, \quad \pi = \pi_{\text{A} \rightarrow \text{C}} = \pi_{\text{C} \rightarrow \text{A}}$$

$F(z, t)$ **rational function**

$$\mathbf{P}(b \in S_n(1) \mid b \notin S_n(0)) = p_n \times \mathbf{P}(\text{no } b \text{ in } S_n(0)) = [z^n] \left. \frac{\partial F(z, t)}{\partial t} \right|_{t=1}$$

$$\mathbf{P}(\text{no } b \text{ in } S_n(0)) = [z^n] F(z, 1)$$

Formal Languages Approach - (N 2013)

(Assuming **a single mutation**)

$b = AAAAA$

$S(0) = XXXX...XXXAAAAAAXAAAAAXXXX.....XXX$

$S(1) = XXXX...XXXAAAAAAXXXX.....XXX$

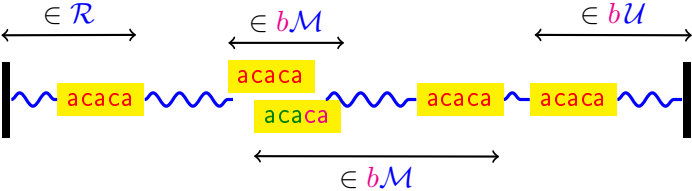
= **XX...** – short clump of **AAAAA** – **...XXX**

- ▶ **length of short clump** of b in $S(1)$ **must be less than** $2 \times |b| - 1$,
- ▶ else **there is at least one occurrence** of b in $S(0)$
- ▶ **no occurrences** of b in the **XXX...XXX**

if b **without self-overlap**, **short clump**= b

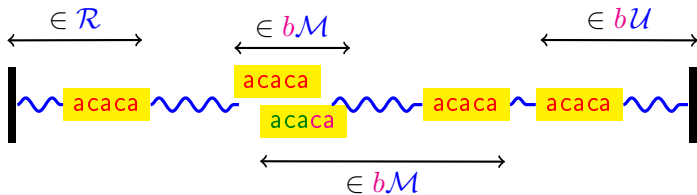
Guibas-Odlyzko decomposition - occurrences of a word b

$b = acaca$



Guibas-Odlyzko decomposition - occurrences of a word b

$b = \text{acaca}$

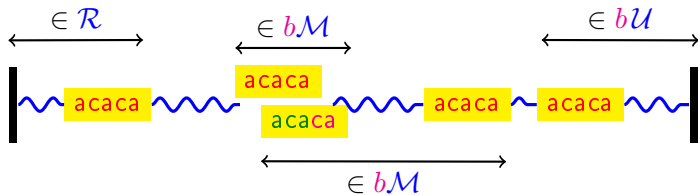


► Right \mathcal{R} : = $\{ w = u.b \text{ et } \nexists r, s, w = r.b.s \}$

$aaaaaacaca \in \mathcal{R}, \quad cccccacacaca \notin \mathcal{R}$

Guibas-Odlyzko decomposition - occurrences of a word b

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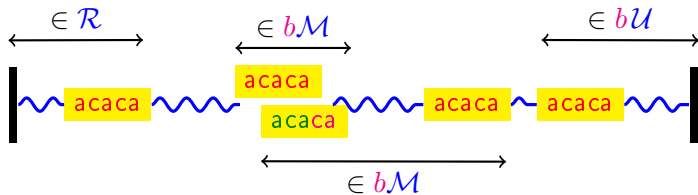
$aaaaaacaca \in \mathcal{R}, \quad cccccacacaca \notin \mathcal{R}$

► Minimal \mathcal{M} : = $\{ w, \quad b.w = u.b \text{ et } \nexists r, s, b.w = r.b.s \}$

$acaca$
 $aaaaacaca \in \mathcal{M} \quad ccaca$
 $caccccccccacaca \notin \mathcal{M} \quad ccaca$
 $ca \in \mathcal{M}$

Guibas-Odlyzko decomposition - occurrences of a word b

$b = \text{acaca}$



- ▶ Right \mathcal{R} : = $\{ w = u.b \text{ et } \nexists r, s, w = r.b.s \}$

$aaaaa\text{acaca} \in \mathcal{R}, \quad ccccc\text{acacaca} \notin \mathcal{R}$

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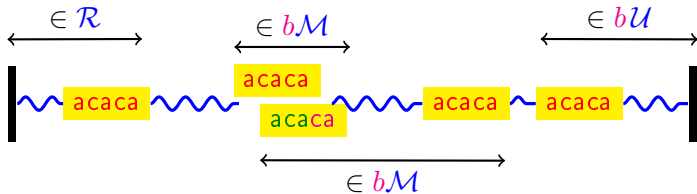
$\text{acaca}aaaa\text{acaca} \in \mathcal{M} \quad ccaca\text{ccccccccacaca} \notin \mathcal{M} \quad ccaca\text{ca} \in \mathcal{M}$

- ▶ Ultimate \mathcal{U} : = $\{ w, \quad \nexists r, s, b.w = r.b.s \}$

$\text{acaca}a\text{accccccccc} \in \mathcal{U} \quad ccaca\text{cccccccc} \notin \mathcal{U}$

Guibas-Odlyzko decomposition - occurrences of a word b

$b = \text{acaca}$



- ▶ Right \mathcal{R} : $= \{ w = u.b \text{ et } \nexists r, s, w = r.b.s \}$

$aaaaaacaca \in \mathcal{R}, \quad cccccacacaca \notin \mathcal{R}$

- ▶ Minimal \mathcal{M} : $= \{ w, \quad b.w = u.b \text{ et } \nexists r, s, b.w = r.b.s \}$

$acaca \quad aaaaaacaca \in \mathcal{M} \quad ccaca \quad cccccccccacaca \notin \mathcal{M} \quad ccaca \quad ca \in \mathcal{M}$

- ▶ Ultimate \mathcal{U} : $= \{ w, \quad \nexists r, s, b.w = r.b.s \}$

$acaca \quad aaccccccccc \in \mathcal{U} \quad ccaca \quad cccccccc \notin \mathcal{U}$

- ▶ Zero $\mathcal{Z} := \mathcal{A}^* - \mathcal{A}^*.b.\mathcal{A}^* = \{ w, \quad \nexists r, s, w = r.b.s \}$

Régnier-Szpankowski Equations (see Lothaire)

- ▶ $A^* = U + MA^*$
- ▶ $A^*b = \mathcal{R}.C + \mathcal{R}.A^*.b$
- ▶ $M^+ = A^*.b + C - \epsilon$
- ▶ $\mathcal{Z}.\sigma = \mathcal{R} + \mathcal{Z} - \epsilon$

Generating Functions of the Languages

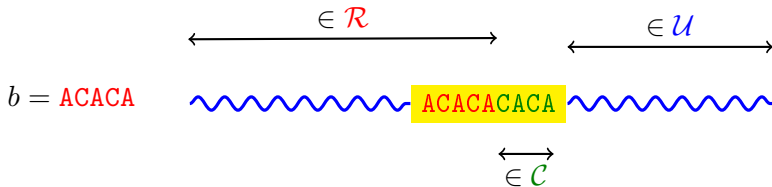
$$\left. \begin{aligned} R(z) &= \frac{\mathbf{P}(b)z^{|b|}}{D(z)}, & M(z) &= 1 - \frac{1-z}{D(z)}, \\ U(z) &= \frac{1}{D(z)}, & Z(z) &= \frac{C(z)}{D(z)}, \end{aligned} \right\} \text{with } D(z) = (1-z)C(z) + \mathbf{P}(b)z^{|b|},$$

\mathcal{C} autocorrelation set of the word b

$$\mathcal{C} = \{w; \quad b.w = u.b, \quad 0 \leq |w| < |b|\}$$

$$C(z) = \sum_{w \in \mathcal{C}} \mathbf{P}(w)z^{|w|}$$

What do we need in $S_n(1)$?

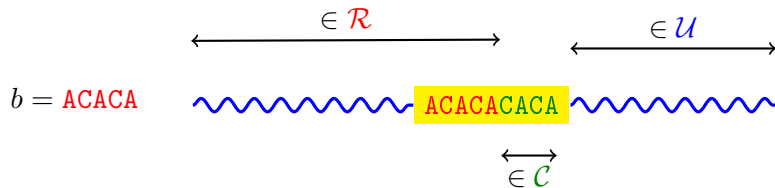


But **not any position** of the clump **can mutate**

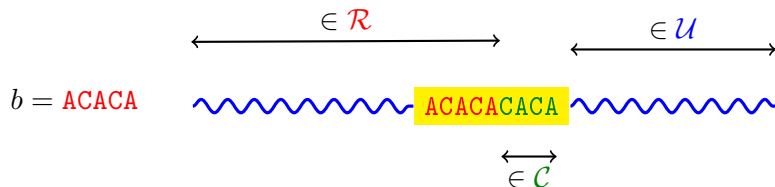
ACACACACA	ACACACA	ACACA
NNNNYNNNN	NNYYYNN	YYYYY

- ▶ **to avoid an occurrence** of b in $S(0)$
- ▶ if the short clump is $b.c$ with $c \in \mathcal{C}$
- ▶ **only $t = |b| - |c|$ positions can mutate**
- ▶ these positions are the t **last positions** of b

The right generating function



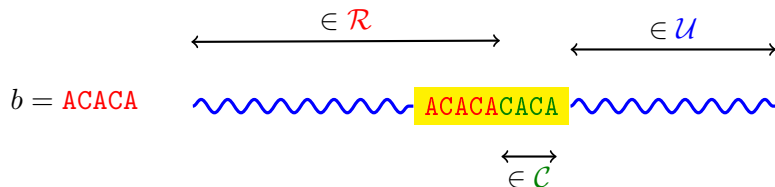
The right generating function



- ▶ Gen.Fun. $F(z)$ of sequences with one short clump

$$F(z) = R(z) \times \sum_{c \in \mathcal{C}} \mathbf{P}(c) z^{|c|} \times U(z) = \sum_{c \in \mathcal{C}} \frac{\mathbf{P}(b.c) z^{|b.c|}}{D^2(z)}$$

The right generating function



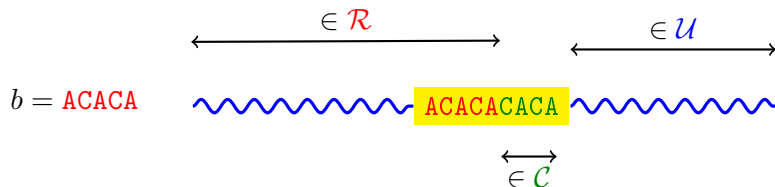
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- ▶ Gen.Fun $Z(z)$ of sequences with no occurrences of b

$$Z(z) = \frac{C(z)}{D(z)}$$

The right generating function



- ▶ Gen.Fun. $F(z)$ of sequences with one short clump

$$F(z) = R(z) \times \sum_{c \in \mathcal{C}} \mathbf{P}(c) z^{|c|} \times U(z) = \sum_{c \in \mathcal{C}} \frac{\mathbf{P}(b.c) z^{|b.c|}}{D^2(z)}$$

- ▶ Gen.Fun $Z(z)$ of sequences with no occurrences of b

$$Z(z) = \frac{C(z)}{D(z)}$$

- ▶ $D(z) = (1 - z)C(z) + \mathbf{P}(b)z^{|b|}$
- ▶ $F(z)$ and $Z(z)$ have the **same dominant singularity** ω

Asymptotics of q_n (approximation of p_n)

$$q_n = \frac{[z^n]F(z)}{[z]Z(z)} \quad (\mathbf{P}(\epsilon) = 1)$$

ω dominant singularity of $D(z)$

$$q_n = \frac{\mathbf{P}(b)}{C(\omega)D'(\omega)} \times \sum_{c \in \mathcal{C}} (|b| - |c|) \mathbf{P}(c) \omega^{|b.c|} \times \sum_{\substack{\beta \in \{b_{[|c|+1]}, \dots, b_{[|b|]}\} \\ \alpha \neq \beta}} \frac{\mathbf{P}(\alpha)}{\mathbf{P}(\beta)} \times \pi_{\alpha \rightarrow \beta} \times \left((n - |b.c| + 1) \omega^{-1} + \frac{D''(\omega)}{D'(\omega)} \right) + o(\mathbf{P}(b)).$$

An even more approximated result

$$D(z) = (1 - z)C(z) + \mathbf{P}(b)z^{|b|}$$

by bootstrapping $\omega \approx 1 + \frac{\mathbf{P}(b)}{C(1) + |b|\mathbf{P}(b)} \approx 1$

Using $\omega \approx 1$ gives

$$\begin{aligned} \mathfrak{q}_n^{(\text{approx})} = & \frac{\mathbf{P}(b)}{C^2(1)} \times \sum_{c \in \mathcal{C}} (|b| - |c|) \mathbf{P}(c) (n - |b \cdot c| + 1) \\ & \times \sum_{\substack{\beta \in \{b_{\lfloor |c|+1} \rfloor}, \dots, b_{\lfloor |b|} \rfloor \\ \alpha \neq \beta}} \frac{\mathbf{P}(\alpha)}{\mathbf{P}(\beta)} \times \pi_{\alpha \rightarrow \beta} \end{aligned}$$

Theorem [N 2013]. The conditioned probability \mathbf{p}_n that a random sequence of length n that does not contain a k -mer b at time 0 evolves at time 1 to a random sequence that contains this k -mer verifies

$$\mathbf{p}_n = \mathbf{q}_n \times (1 + \mathcal{O}(n\psi)) + \mathcal{O}(n^2\psi^2)$$

where

$$\begin{aligned} \mathbf{q}_n &= \frac{\mathbf{P}(b)}{C(\omega)D'(\omega)} \\ &\times \sum_{c \in \mathcal{C}} (|b| - |c|) \mathbf{P}(c) \omega^{|b \cdot c|} \times \sum_{\substack{\beta \in \{b_{[|c|+1]}, \dots, b_{[|b|]}\} \\ \alpha \neq \beta}} \frac{\mathbf{P}(\alpha)}{\mathbf{P}(\beta)} \times \pi_{\alpha \rightarrow \beta} \\ &\times \left((n - |b \cdot c| + 1) \omega^{-1} + \frac{D''(\omega)}{D'(\omega)} \right) + o(\mathbf{P}(b)). \end{aligned}$$

$$\psi = \frac{\max_{\alpha, \beta \in \mathcal{A}; \alpha \neq \beta} p_{\alpha \rightarrow \beta}}{\min_{\alpha \in \mathcal{A}} p_{\alpha \rightarrow \alpha}}$$

Numerical validation

$\mathcal{A} = \{A, C, G, T\}$ - uniform Bernoulli model for $S(0)$.

	$b = \text{AAAAA}$ and		for $\alpha \neq \beta$, $p_{\alpha \rightarrow \beta} = 10^{-8}$	
Length n	$\mathfrak{p}_n \times 10^6$	$\mathfrak{h}_n \times 10^6$	$\mathfrak{q}_n \times 10^6$	$\mathfrak{q}_n^{(\text{approx})} \times 10^6$
10000	1.03335528	1.03335588	1.03335587	1.02703244
100000	10.3368481	10.3369021	10.3369021	10.2742439
1000000	1033.19278	1033.72699	1033.72698	1027.46750

- ▶ \mathfrak{p}_n - Exact result by automata (Behrens-Nicaud-N 2012)
- ▶ Heuristic of a single mutation
 - ▶ \mathfrak{h}_n clumps of neighbors at distance 1 of b in $S_n(0)$ (N 2012)
 - ▶ $\mathfrak{q}_n, \mathfrak{q}_n^{(\text{approx})}$ short clump approach on $S_n(1)$ (N 2103)



Bruno raised **vigorous** and **justified critics** to a previous version of this talk.



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So, if you did not like it, **please complain to Bruno**