

Diagonal asymptotics for products of combinatorial classes

Or: the diagonal method is still not very good

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- ▶ No Maple package, but there is now a reasonable implementation in Sage (available at Alex's website). Needs some algorithmic speedups. Any volunteers?
- ▶ In 2012, I saw that the word has not yet spread far enough. Multivariate methods are more general, conceptually simpler, and, I claim, computationally superior.

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- ▶ The answer can be given explicitly in this case:

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- ▶ Suppose we replace “two” by d , \mathbb{N} by other combinatorial classes, allow different n for different compositions, ... ?

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- ▶ Banderier & Hitczenko 2012: generalize from 2 to d compositions, different restriction S for each one. Some explicit formulae and asymptotics.

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- ▶ Use the symbolic method. Let $F(\mathbf{x}, \mathbf{y}) = \sum a_{\mathbf{n}} \mathbf{x}^{\mathbf{n}} \mathbf{y}^{\mathbf{k}}$ be the $2d$ -variate generating function, where \mathbf{x} marks size and \mathbf{y} marks number of components. Here $F(\mathbf{x}, \mathbf{y})$ factors as $\prod_{i=1}^d F_i(x_i, y_i)$.

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- ▶ The number of d -tuples of objects with the same number of components is $[\mathbf{x}^{\mathbf{n}}] \text{diag}_{\mathbf{y}} F(\mathbf{x}, \mathbf{1})$. In particular for the simplest case where all $n_i = n$,

$$[\mathbf{x}^{n\mathbf{1}}] \text{diag}_{\mathbf{y}} F(\mathbf{x}, \mathbf{1}) = \sum_{k \geq 0} (a_{nk})^d =: b_n.$$

Aside: exact solutions

- ▶ When $d = 2$, we have a good chance of finding an exact solution. For Dyck walks

$$\sum_{\substack{0 \leq k \leq n \\ 2 \mid (n-k)}} \left[\frac{k+1}{n+1} \binom{n+1}{\frac{n-k}{2}} \right]^2 = \frac{1}{n+1} \binom{2n}{n}.$$

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More generally, when (a_{nk}) is a **Riordan array**, namely the case $F_i(x, y) = \phi(x)/(1 - yv(x))$, we discover new identities of this type that are not in OEIS.

- ▶ When $d \geq 3$, exact solutions are rare. For example, $b_n = \sum_k \binom{n}{k}^3$ is known not to have an algebraic generating function.

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- ▶ It seems that the work needed is enormous even for rather modest-looking problems. For example, the defining linear differential equation for $\sum_k \binom{n-k}{k}^5$ has order 6 with polynomial coefficients of degree 38. Banderier and Hitczenko report: “Current state of the art algorithms will take more than one day for $d = 6$, and gigabytes of memory”

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- ▶ How to do it for general d ? Also, the diagonal method does not yield asymptotics that are uniform in the slope of the diagonal; performance away from the main diagonal is bad.

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- ▶ Consider the random variable X_n whose PGF is $\sum_k a_{nk} y^k / \sum_k a_{nk}$, mean μ_n , variance σ_n^2 . If $(X_n - \sigma_n) / \mu_n$ converges to a continuous limit law with density g , then

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- ▶ In the Gaussian case, K is explicitly computable.
- ▶ As usual, such methods say nothing about higher order terms, or when there is not a continuous limit. Still, this approach is a useful complement to the above methods.

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 - ▶ Asymptotics of algebraic functions via lifting to a rational function in higher dimension (resolution of singularities).
 - ▶ We want numerical approximations for smaller values of n .
- ▶ This topic was the subject of two papers with Alex Raichev. For example, our 2nd order approximation for $\sum_{k=0}^n \binom{n}{5}$ even for $n = 8$ has relative error only 0.5%, but 10% for 1st order.

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- ▶ For more, see the book (next talk!).

General asymptotic formula (supercritical Riordan case)

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- ▶ Suppose $F_i(x, y) = \phi(x)/(1 - yv(x))$ and ϕ has radius of convergence large enough. Let $c > 0$ solve $v(c) = 1$. Then

$$b_{n\mathbf{1}} \sim c^{-dn} n^{-d/2} \sum_l c_l n^{-l} \quad \text{where } c_l \text{ is explicitly computable.}$$

In particular

$$c_0 = \frac{\phi(c)^d}{\sqrt{d}\mu_v(c) \left[2\pi \frac{\sigma_v^2(c)}{\mu_v(c)} \right]^{\frac{d-1}{2}}}.$$

Examples



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$$\sum_{k=0}^n \binom{n}{k}^6 \sim 64^n \left(\frac{4\sqrt{3}}{3(\pi n)^{\frac{5}{2}}} - \frac{25\sqrt{3}}{9\pi^{\frac{5}{2}} n^{\frac{7}{2}}} \right)$$

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- ▶ Once the diagonal GF is found, the asymptotic extraction is quicker, since it is a univariate problem. The multivariate method typically requires solving systems of algebraic equations.
- ▶ I suggest a serious theoretical and experimental comparison of the performance of these methods. If done experimentally, we need to implement the methods equally. I know which one I would bet on to win!